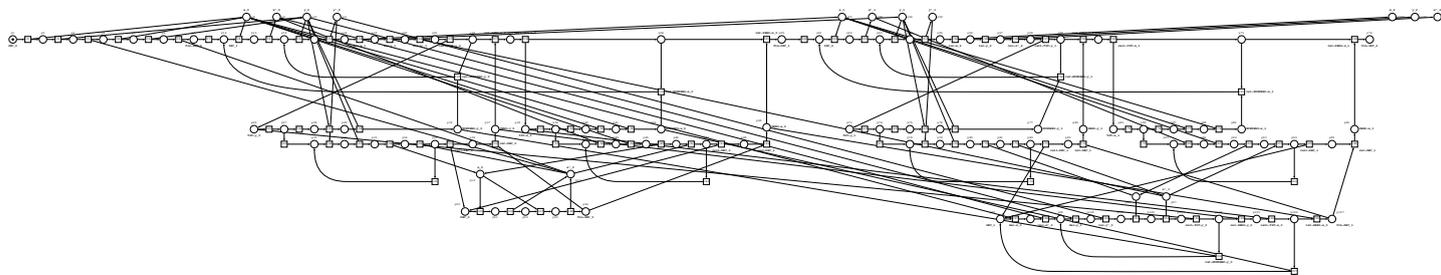


This form is a summary description of the model entitled “DoubleExponent” proposed for the Model Checking Contest @ Petri Nets. Models can be given in several instances parameterized by scaling parameters. Colored nets can be accompanied by one or many equivalent, unfolded P/T nets. Models are given together with property files (possibly, one per model instance) giving a set of properties to be checked on the model.

Description

The model illustrates difference between weak, strong, and exact Petri net computers [1]. The generated for parameter k net [2] counts 2^{2^k} as a strong computer after Richard J. Lipton [3]. Well-structured net programs of Javier Esparza [4] have been encoded by Petri nets. The following marking belongs to the reachability space: $finINC_k = 1$; $x_0 = 2$, $y_0 = 2$, $s'_0 = 2$; $x_1 = 4$, $y_1 = 4$, $s'_1 = 4$; ...; $x_k = 2^{2^k}$, $y_k = 2^{2^k}$, $s'_k = 2^{2^k}$. It is the only marking with $finINC_k = 1$. For a given k , the resulting net is composed of a single $Block_0$ and repeated $(k - 1)$ times $Block_i$. $Block_i$ includes subnets INC_{i+1} , $TSTy_i$, $TSTx_i$, DEC_i and variables (places) x_i , x'_i , y_i , y'_i , s_i , s'_i . Connections of neighboring blocks reflect initialization of the next variables x_{i+1} , y_{i+1} , s'_{i+1} and recursive calls of the previous block's $TSTy_{i-1}$, $TSTx_{i-1}$. Models are generated by program [2] which implements technique described in [5].



Graphical representation for $k = 2$

References

- [1] Dmitry Zaitsev, Some Remarks on Petri Net Computers: Weak, Exact, and Strong. Petri Net Newsletter, Volume 85, December 2016, 3-7.
- [2] Dmitry Zaitsev, Generator of Petri nets which count double exponent 2^{2^k} after R.J.Lipton and J.Esparza constructs. <https://github.com/dazeorgacm/depn>
- [3] Richard J. Lipton, The Reachability Problem Requires Exponential Space, Technical Report 63, Yale University, 1976. <http://www.cs.yale.edu/publications/techreports/tr63.pdf>
- [4] Javier Esparza, Decidability and Complexity of Petri Net Problems - An Introduction. LNCS 1491, 1996, 374-428. <http://www7.in.tum.de/um/bibdb/esparza/course.pdf>
- [5] Dmitry A. Zaitsev, Ivan D. Zaitsev and Tatiana R. Shmeleva. Infinite Petri Nets: Part 2, Modeling Triangular, Hexagonal, Hypercube and Hypertorus Structures, Complex Systems, 26(4), 2017, 341-371.

Scaling parameter

Parameter name	Parameter description	Chosen parameter values
k	k – power of double exponent 2^{2^k} to compute	1,2,3,4,10,20,100,200

Size of the model

Parameter	Number of places	Number of transitions	Number of arcs
k	$P = 53(k - 1) + 57$	$T = 50(k - 1) + 48$	$A = 141(k - 1) + 135$
$k = 1$	57	48	135
$k = 2$	110	98	276
$k = 3$	163	148	417
$k = 4$	216	198	558
$k = 10$	534	498	1404
$k = 20$	1064	998	2814
$k = 100$	5304	4998	14094
$k = 200$	10604	9998	28194

Structural properties

- ordinary** — all arcs have multiplicity one yes
- simple free choice** — all transitions sharing a common input place have no other input place no ^(a)
- extended free choice** — all transitions sharing a common input place have the same input places no ^(b)
- state machine** — every transition has exactly one input place and exactly one output place no ^(c)
- marked graph** — every place has exactly one input transition and exactly one output transition no ^(d)
- connected** — there is an undirected path between every two nodes (places or transitions) yes ^(e)
- strongly connected** — there is a directed path between every two nodes (places or transitions) no ^(f)
- source place(s)** — one or more places have no input transitions yes ^(g)
- sink place(s)** — one or more places have no output transitions yes ^(h)
- source transition(s)** — one or more transitions have no input places no ⁽ⁱ⁾
- sink transitions(s)** — one or more transitions have no output places no ^(j)
- loop-free** — no transition has an input place that is also an output place yes ^(k)
- conservative** — for each transition, the number of input arcs equals the number of output arcs no ^(l)
- subconservative** — for each transition, the number of input arcs equals or exceeds the number of output arcs no ^(m)
- nested units** — places are structured into hierarchically nested sequential units ⁽ⁿ⁾ no

Behavioural properties

- safe** — in every reachable marking, there is no more than one token on a place no ^(o)
- dead place(s)** — one or more places have no token in any reachable marking ?
- dead transition(s)** — one or more transitions cannot fire from any reachable marking no
- deadlock** — there exists a reachable marking from which no transition can be fired yes
- reversible** — from every reachable marking, there is a transition path going back to the initial marking no
- live** — for every transition t , from every reachable marking, one can reach a marking in which t can fire no

^(a) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(b) transitions “t10” and “t23” share a common input place “p10”, but only the former transition has input place “p16”.

^(c) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(d) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(e) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(f) from place “p10” one cannot reach place “p1”.

^(g) place “p1” is a source place.

^(h) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

⁽ⁱ⁾ stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(j) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(k) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(l) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

^(m) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

⁽ⁿ⁾ the definition of Nested-Unit Petri Nets (NUPN) is available from <http://mcc.lip6.fr/nupn.php>

^(o) stated by [CÆSAR.BDD](#) version 2.7 on all 8 instances (for $k = 1, 2, 3, 4, 10, 20, 100$, and 200).

Size of the marking graphs

Parameter	Number of reachable markings	Number of transition firings	Max. number of tokens per place	Max. number of tokens per marking
k	?	?	2^{2^k}	$> F, F = 3 \cdot \sum_{i=0}^k 2^{2^i} + 1$ (guess: $F + 2k$)
1	149	148	4	21
2	3708	3707	16	71
3	2,385,072	2,385,071	256	841
4	$\geq 27^{(p)}$?	65,536	≥ 8
10	$\geq 27^{(q)}$?	$\approx 1.797693135 \times 10^{308}$	≥ 8
20	$\geq 27^{(r)}$?	$\sim 10^{314.573}$	≥ 8
100	?	?	$\sim 10^{38 \cdot 10^{28}}$	≥ 8
200	?	?	$\sim 10^{482 \cdot 10^{57}}$	≥ 8

Other properties

The net is a Strong Computer of 2^{2^k} and its size is linear in k that proves exponential complexity in space of Petri net reachability problem.

^(p) stated by [CÆSAR.BDD](#) version 3.3.

^(q) stated by [CÆSAR.BDD](#) version 3.3.

^(r) stated by [CÆSAR.BDD](#) version 3.3.