

This form is a summary description of the model entitled “ARMCACHECoherency” proposed for the Model Checking Contest @ Petri Nets. Models can be given in several instances parameterized by scaling parameters. Colored nets can be accompanied by one or many equivalent, unfolded P/T nets. Models are given together with property files (possibly, one per model instance) giving a set of properties to be checked on the model.

Description

The ACE specification proposed by ARM is becoming a de facto industrial standard for system-level cache coherence in heterogeneous systems on chip.

The present P/T net was derived from a formal specification of ACE developed by STMicroelectronics and INRIA Grenoble/LIG laboratory. This specification is written in LNT (*LOTOS New Technology*), which combines functional languages (to describe data types and user-defined functions operating on typed values) and process calculi (to describe concurrent components that synchronize using rendezvous and communicate via message passing). The LNT specification is 2670-line long and models a system with two ACE masters and one ACE-Lite master, where:

- global constraints are included;
- one ACE master may execute MakeUnique, ReadOnce, and WriteBack transactions;
- the other ACE master executes no transactions;
- the ACE-Lite master executes ReadOnce transactions.

The LNT specification was translated to LOTOS, and then to an interpreted Petri net using the CADP toolbox. Finally, the present P/T net was obtained by stripping out all dataflow-related information (variables, types, assignments, guards, etc.) from the interpreted Petri net, leading to a NUPN (*Nested-Unit Petri Net*) model translated to PNML using the CÆSAR.BDD tool.

References

Abderahman Kriouile and Wendelin Serwe. *Formal Analysis of the ACE Specification for Cache Coherent Systems-on-Chip*. Proceedings of the 18th International Workshop on Formal Methods for Industrial Critical Systems (FMICS'13). Lecture Notes in Computer Science 8187, pp. 108-122, 2013. <http://hal.inria.fr/hal-00858521/en>

ARM. *AMBA AXI and ACE Protocol Specification*, version ARM IHI 0022E, February 2013. <http://infocenter.arm.com/help/topic/com.arm.doc.ih0022e>

Scaling parameter

This model is not parameterized.

Size of the model

number of places:	87
number of transitions:	33676
number of arcs:	246935
number of units:	13
HWB code (<i>height-width-bits</i>):	2-12-37

Structural properties

ordinary — all arcs have multiplicity one	yes	
simple free choice — all transitions sharing a common input place have no other input place	no	(a)
extended free choice — all transitions sharing a common input place have the same input places	no	(b)
state machine — every transition has exactly one input place and exactly one output place	no	(c)
marked graph — every place has exactly one input transition and exactly one output transition	no	(d)
connected — there is an undirected path between every two nodes (places or transitions)	yes	(e)
strongly connected — there is a directed path between every two nodes (places or transitions)	no	(f)
source place(s) — one or more places have no input transitions	yes	(g)
sink place(s) — one or more places have no output transitions	no	(h)
source transition(s) — one or more transitions have no input places	no	(i)
sink transitions(s) — one or more transitions have no output places	no	(j)
loop-free — no transition has an input place that is also an output place	no	(k)
conservative — for each transition, the number of input arcs equals the number of output arcs	no	(l)
subconservative — for each transition, the number of input arcs equals or exceeds the number of output arcs	no	(m)
nested units — places are structured into hierarchically nested sequential units ⁽ⁿ⁾	yes	

Behavioural properties

safe — in every reachable marking, there is no more than one token on a place	yes	(o)
dead place(s) — one or more places have no token in any reachable marking	no	(p)
dead transition(s) — one or more transitions cannot fire from any reachable marking	no	(q)
deadlock — there exists a reachable marking from which no transition can be fired	no	(r)
reversible — from every reachable marking, there is a transition path going back to the initial marking	?	
live — for every transition t , from every reachable marking, one can reach a marking in which t can fire	?	

Size of the marking graph

number of reachable markings:	3.206E+8 ^(s)
number of transition firings:	2.234E+10 ^(t)
max. number of tokens per place:	1 ^(u)
max. number of tokens per marking:	12

(a) 123325 arcs are not simple free choice, e.g., the arc from place 1 (which has 4425 outgoing transitions) to transition 132 (which has 3 input places).

(b) transitions 12232 and 132 share a common input place 1, but only the former transition has input place 8.

(c) 33541 transitions are not of a state machine, e.g., transition 0.

(d) 87 places are not of a marked graph, e.g., place 0.

(e) stated by [CÆSAR.BDD](#) version 1.5.

(f) from place 1 one cannot reach place 0.

(g) place 0 is a source place.

(h) stated by [CÆSAR.BDD](#) version 1.5.

(i) stated by [CÆSAR.BDD](#) version 1.5.

(j) stated by [CÆSAR.BDD](#) version 1.5.

(k) 22262 transitions are not loop free, e.g., transition 6.

(l) transition 0 is not conservative.

(m) transition 0 is not subconservative.

(n) the definition of Nested-Unit Petri Nets (NUPN) is available from <http://mcc.lip6.fr/nupn.php>

(o) safe by construction – stated by the [CÆSAR](#) compiler.

(p) stated by [CÆSAR.BDD](#) version 3.3.

(q) stated by [CÆSAR.BDD](#) version 2.0.

(r) stated by [CÆSAR.BDD](#) version 2.0.

(s) stated by [CÆSAR.BDD](#) version 2.0; confirmed at MCC'2014 by Marcie and PNMC.

(t) computed at MCC'2014 by Marcie.

(u) confirmed at MCC'2014 by GretSPN and Marcie.