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This form is a summary description of the model entitled "JoinFreeModules" proposed for the Model Checking Contest @ Petri Nets. Models can be given in several instances parameterized by scaling parameters. Colored nets can be accompanied by one or many equivalent, unfolded P/T nets. Models are given together with property files (possibly, one per model instance) giving a set of properties to be checked on the model.

Description

Liveness and boundedness are well-known Petri net properties that are fundamental for many real-world applications, including embedded systems, and are hard to check. So as to alleviate this difficulty, particular subclasses are often considered. The weighted join-free class has a limited expressiveness, since it forbids synchronizations. However, even with this strong structural limitation, this class remains hard to analyze and appears as a fundamental module of more complex classes, such as weighted asymmetric-choice nets whose behavior is strongly related to its join-free modules [1].

The model provided is inspired from a paper of T. Hujsa and R. Devillers [1] which investigates the relationship between liveness and structural boundedness (meaning boundedness for every initial marking) in some subclasses of weighted Petri nets, notably join-free nets. The paper also studies the monotonicity of liveness, meaning its preservation upon any increase of the live marking considered.

The model is based on one parameter k, which denotes:

- the number of join-free modules,
- the initial number of tokens in each buffer (place) $b_1, \ldots, b_k$
- the amount of tokens read and written (as a self-loop) through one firing of t (outside the modules) in each $b_1, \ldots, b_k$.

Each module N_i, $i \in \{1, \ldots, k\}$, is a structurally live and bounded join-free net containing the buffer b_i. A single transition t checks the existence of i tokens in each buffer b_i. Each N_i can represent a sub-program with asynchronous iterations to be executed concurrently on different processors. Each new iteration in N_i reads some data items (tokens) in its buffer b_i.

The sub-program t reads k data items in each buffer b_i, computes a function on them and updates the buffer b with the result. The purpose of this operation is to analyse the current progress of each module and to gather this updated information in b. Due to this synchronisation t, the global system is not join-free.

Each join-free module N_i is live for $k = 3$, becomes non-live for $k = 4, 5, 6$ (proving that liveness is not monotonic in the structurally bounded join-free class) and becomes live again for all $k \geq 7$ (this stems partly from a result of [2]). Each N_i is live and reversible for all $k \geq 8$, implying that the global system is live and reversible for every $k \geq 8$, at least."

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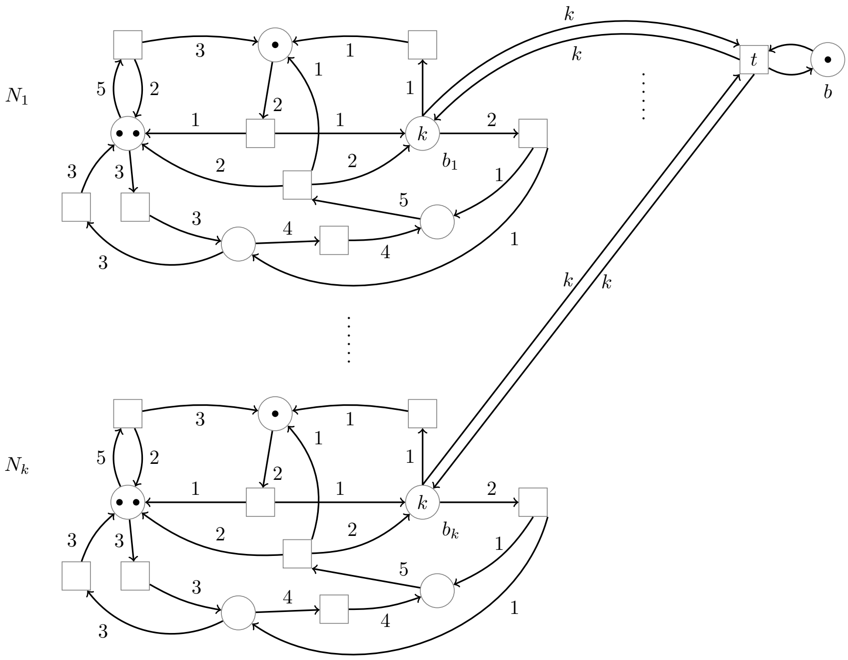

For each $i \in \{1, \dots, k\}$, the subnet N_i contains a place b_i with k initial tokens. The transition t can fire when at least k tokens are present in each b_i.

References

[1] Thomas Hujsa and Raymond Devillers. *On Liveness and Deadlockability in Subclasses of Weighted Petri Nets.* Proceedings of the 38th International Conference on Application and Theory of Petri Nets and Concurrency (Petri Nets'17), 2017.

[2] Jean-Marc Delosme, Thomas Hujsa, and Alix Munier-Kordon. *Polynomial sufficient conditions of well-behavedness for weighted Join-Free and Choice-Free systems.* In Proceedings of the 13th International Conference on Application of Concurrency to System Design (ACSD'13). pages 90–99, 2013.

Scaling parameter

Parameter name	**Parameter description**	**Chosen parameter values**
k	The number of initial tokens in each buffer b_i (in the module N_i), the number of tokens consumed and produced by each firing of t in the same buffers, and the number of modules	3, 4, 5, 10, 20, 50, 100, 200, 500, 1000, 2000, 5000

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Size of the model

Parameter	Number of places	Number of transitions	Number of arcs
k	$5k+1$	$8k+1$	$23k+2$
$k=3$	16	25	71
$k=4$	21	33	94
$k=5$	26	41	117
$k=10$	51	81	232
$k=20$	101	161	462
$k=50$	251	401	1 152
$k=100$	501	801	2 302
$k=200$	1 001	1 601	4 602
$k=500$	2 501	4 001	11 502
$k=1000$	5 001	8 001	23 002
$k=2000$	10 001	16 001	46 002
$k=5000$	25 001	40 001	115 002

Structural properties

ordinary — *all arcs have multiplicity one* .... ✗
simple free choice — *all transitions sharing a common input place have no other input place* .... ✗ (a)
extended free choice — *all transitions sharing a common input place have the same input places* .... ✗ (b)
state machine — *every transition has exactly one input place and exactly one output place* .... ✗ (c)
marked graph — *every place has exactly one input transition and exactly one output transition* .... ✗ (d)
connected — *there is an undirected path between every two nodes (places or transitions)* .... ✓ (e)
strongly connected — *there is a directed path between every two nodes (places or transitions)* .... ✓ (f)
source place(s) — *one or more places have no input transitions* .... ✗ (g)
sink place(s) — *one or more places have no output transitions* .... ✗ (h)
source transition(s) — *one or more transitions have no input places* .... ✗ (i)
sink transitions(s) — *one or more transitions have no output places* .... ✗ (j)
loop-free — *no transition has an input place that is also an output place* .... ✗ (k)
conservative — *for each transition, the number of input arcs equals the number of output arcs* .... ✗ (l)
subconservative — *for each transition, the number of input arcs equals or exceeds the number of output arcs* .... ✓ (m)
nested units — *places are structured into hierarchically nested sequential units* (n) .... ✗

Behavioural properties

safe — *in every reachable marking, there is no more than one token on a place* .... ✗ (o)
deadlock — *there exists a reachable marking from which no transition can be fired* .... ?
reversible — *from every reachable marking, there is a transition path going back to the initial marking* .... ?
quasi-live — *for every transition t, there exists a reachable marking in which t can fire* .... ?

(a) the net is not ordinary.
(b) the net is not ordinary.
(c) the net is not ordinary.
(d) the net is not ordinary.
(e) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(f) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(g) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(h) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(i) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(j) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(k) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).
(l) stated by PNML2NUPN 1.5.3 on all 12 instances (see all aforementioned parameter values).
(m) stated by PNML2NUPN 1.5.3 on all 12 instances (see all aforementioned parameter values).
(n) the definition of Nested-Unit Petri Nets (NUPN) is available from `http://mcc.lip6.fr/nupn.php`
(o) stated by CÆSAR.BDD version 2.7 on all 12 instances (see all aforementioned parameter values).

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live — *for every transition t, from every reachable marking, one can reach a marking in which t can fire* ....?

Size of the marking graphs

Parameter	Number of reachable markings	Number of transition firings	Max. number of tokens per place	Max. number of tokens per marking
For any k:	–	–	$\leq k+3$	$k \cdot (k+3)+1$
$k=3$	35 937	225 450	≤ 6	19 [p]
$k=4$	14 776 336	138 230 321	≤ 7	29 [q]
$k=5$	?	?	≤ 8	41 [r]
$k=10$	?	?	≤ 13	131 [s]
$k=20$	?	?	≤ 23	461 [t]
$k=50$	?	?	≤ 53	2 651 [u]
$k=100$	?	?	≤ 103	10 301 [v]
$k=200$	?	?	≤ 203	40 601 [w]
$k=500$	?	?	≤ 503	251 501 [x]
$k=1000$	?	?	$\leq 1\,003$	1 003 001 [y]
$k=2000$	?	?	$\leq 2\,003$	4 006 001 [z]
$k=5000$	?	?	$\leq 5\,003$	2 5015 001 [aa]

Other properties

P — and T – invariants exist and cover the net. Denoting by L the liveness property and by R the reversibility property, there exist markings satisfying:

- L and R
- L and not R
- not L and R
- not L and not R

For the value $k=4$, a deadlock is reachable.

[p] number of initial tokens, because the net is sub-conservative.
[q] number of initial tokens, because the net is sub-conservative.
[r] number of initial tokens, because the net is sub-conservative.
[s] number of initial tokens, because the net is sub-conservative.
[t] number of initial tokens, because the net is sub-conservative.
[u] number of initial tokens, because the net is sub-conservative.
[v] number of initial tokens, because the net is sub-conservative.
[w] number of initial tokens, because the net is sub-conservative.
[x] number of initial tokens, because the net is sub-conservative.
[y] number of initial tokens, because the net is sub-conservative.
[z] number of initial tokens, because the net is sub-conservative.
[aa] number of initial tokens, because the net is sub-conservative.