

This form is a summary description of the model entitled “ClientsAndServers” proposed for the Model Checking Contest @ Petri Nets. Models can be given in several instances parameterized by scaling parameters. Colored nets can be accompanied by one or many equivalent, unfolded P/T nets. Models are given together with property files (possibly, one per model instance) giving a set of properties to be checked on the model.

Description

Let be a set C of client, a set S of servers, a set M of managers and a pool of U resource units. A client may send a request a resource to the server set. Any server may transmit the request to the manager set and waits for a grant. Any manager may allocate a resource unit from the pool and returns a grant. A waiting server transmits this grant to the client. After using the resource the client frees it by sending a message to the server set. Any server may notify the manager set and waits for an acknowledge that it transmits it to the client.

An idle client (place **Ci**) sends (transition **csR**) a client request (place **CR**) to the server set and waits (place **CwG**) until it receives (transition **crG**) a grant (place **CG**) from the server allowing access to a resource unit.

- An idle server (place **Si**) may receive (transition **srR**) a request from a client, treats it (place **StR**) and sends (transition **ssR**) a server request (place **SR**) to the manager set. Then, it waits (place **SwG**) for a grant (place **SG**) from a manager.
- An idle manager (place **Mi**) may receive (transition **mrR**) a server request, it waits (place **MwU**) for a free resource unit (place **Uf**) that it allocates (transition **maU**). Then, it prepares a grant (place **MpG**) that it sends (transition **msG**) to the server and becomes again idle.
- When a waiting server receives (transition **srG**) a grant (place **SG**) from the manager, it prepares (place **SpG**) a grant (place **CG**) for the client, sends it (transition **ssG**), and becomes again in state idle.

When a waiting client receives (transition **crG**) a grant (place **CG**) from a server, it becomes busy (place **Cb**) using the resource unit. After finishing it sends (transition **csF**) a free message (place **CF**) to a server and waits (place **CwA**) for an acknowledge.

- When an idle server (place **Si**) receives (transition **srF**) a free message (place **CF**) from a client, it treats it (place **StF**) and sends (transition **ssF**) a server free message (place **SF**) to the manager set. Then, it waits (place **Sw**) for an acknowledge (place **SA**).
- When an idle manager receives (transition **mrF**) a free message (place **SF**) from a server, it treats it (place **MtF**) and releases (transition **mfree**) the resource. Then, it prepares (place **MpA**) an acknowledge (place **SA**) for the server, sends it (transition **msA**), and becomes again in state idle.
- When a waiting server (place **SwA**) receives (transition **srA**) an acknowledge (place **SA**) from the manager. It prepares (place **SpA**) an acknowledge (place **CA**) for the client, sends it (transition **ssA**), and becomes again in state idle.

When a waiting client (place **CwA**) receives (transition **crA**) this acknowledge, it becomes again in state idle.

In its most general form, the model has four parameters: the number $|C|$ of clients, the number $|S|$ of servers, the number $|M|$ of managers, and the number $|U|$ of resources units.

We consider here a simplified form, in which the model is parameterized by two natural numbers N and P (with $N > P$) such that

- the number of clients is set to $C = 8N$
- the number of servers is set to $S = 2N$
- the number of managers is set to $M = 3N - 3P$
- the number of resource units is set to $U = 4N + 8P$.

Structural properties

ordinary — all arcs have multiplicity one	✓
simple free choice — all transitions sharing a common input place have no other input place	✗ (a)
extended free choice — all transitions sharing a common input place have the same input places	✗ (b)
state machine — every transition has exactly one input place and exactly one output place	✗ (c)
marked graph — every place has exactly one input transition and exactly one output transition	✗ (d)
connected — there is an undirected path between every two nodes (places or transitions)	✓ (e)
strongly connected — there is a directed path between every two nodes (places or transitions)	✓ (f)
source place(s) — one or more places have no input transitions	✗ (g)
sink place(s) — one or more places have no output transitions	✗ (h)
source transition(s) — one or more transitions have no input places	✗ (i)
sink transitions(s) — one or more transitions have no output places	✗ (j)
loop-free — no transition has an input place that is also an output place	✓ (k)
conservative — for each transition, the number of input arcs equals the number of output arcs	✗ (l)
subconservative — for each transition, the number of input arcs equals or exceeds the number of output arcs	✗ (m)
nested units — places are structured into hierarchically nested sequential units ⁽ⁿ⁾	✗

Behavioural properties

safe — in every reachable marking, there is no more than one token on a place	✗ (o)
deadlock — there exists a reachable marking from which no transition can be fired	? (p)
reversible — from every reachable marking, there is a transition path going back to the initial marking	?
quasi-live — for every transition t , there exists a reachable marking in which t can fire	✓ (q)
live — for every transition t , from every reachable marking, one can reach a marking in which t can fire	?

(a) 4 arcs are not simple free choice, e.g., the arc from place “Mi” (which has 2 outgoing transitions) to transition “mrF” (which has 2 input places).

(b) transitions “mrF” and “mrR” share a common input place “Mi”, but only the former transition has input place “SF”.

(c) 18 transitions are not of a state machine, e.g., transition “ssF”.

(d) 2 places are not of a marked graph, e.g., place “Mi”.

(e) confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(f) stated by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(g) confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(h) confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(i) confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(j) confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(k) stated by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(l) 18 transitions are not conservative, e.g., transition “ssF”.

(m) 9 transitions are not subconservative, e.g., transition “ssF”.

(n) the definition of Nested-Unit Petri Nets (NUPN) is available from <http://mcc.lip6.fr/nupn.php>

(o) in the initial marking, some places have several tokens; confirmed by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

(p) depends on the initial marking.

(q) stated by [CÆSAR.BDD](#) version 2.7 on all 20 instances (see aforementioned parameter values).

Size of the marking graphs

Parameter	Number of reach- able markings	Number of tran- sition firings	Max. number of tokens per place	Max. number of tokens per marking
$N = 1, P = 0$	27 576 ^(r)	113 316 ^(s)	?	≥ 17 ^(t)
$N = 2, P = 0$	7 081 638 ^(u)	44 030 250 ^(v)	?	≥ 34 ^(w)
$N = 2, P = 1$	12 462 173 ^(x)	77 859 168 ^(y)	?	≥ 39 ^(z)
$N = 5, P = 0$	$> 39\,919\,315$ ^(aa)	$> 240\,893\,998$ ^(ab)	?	≥ 85 ^(ac)
$N = 5, P = 1$	$> 24\,676\,885$ ^(ad)	$> 161\,268\,212$ ^(ae)	?	≥ 90 ^(af)
$N = 10, P = 0$?	?	?	≥ 170 ^(ag)
$N = 10, P = 1$?	?	?	≥ 175 ^(ah)
$N = 10, P = 2$?	?	?	≥ 180 ^(ai)
$N = 20, P = 0$?	?	?	≥ 340 ^(aj)
$N = 20, P = 1$?	?	?	≥ 345 ^(ak)
$N = 20, P = 2$?	?	?	≥ 350 ^(al)
$N = 20, P = 3$?	?	?	≥ 355 ^(am)
$N = 20, P = 4$?	?	?	≥ 360 ^(an)
$N = 50, P = 0$?	?	?	≥ 850 ^(ao)
$N = 100, P = 0$?	?	?	≥ 1700 ^(ap)
$N = 200, P = 0$?	?	?	≥ 3400 ^(aq)
$N = 500, P = 0$?	?	?	≥ 8500 ^(ar)
$N = 1000, P = 0$?	?	?	≥ 17000 ^(as)
$N = 2000, P = 0$?	?	?	≥ 34000 ^(at)
$N = 5000, P = 0$?	?	?	≥ 85000 ^(au)

Other properties

There are deadlocks iff $|C| \geq |S| + |U|$ and $|C| \geq |M| + |U|$.

The types of deadlock also depend whether $|S| \geq |M|$ and whether $|S| \geq |M| + |U|$.

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- ^(r) stated by prod in April 2017.
 - ^(s) stated by prod in April 2017.
 - ^(t) lower bound given by the number of initial tokens.
 - ^(u) stated by prod in April 2017.
 - ^(v) stated by prod in April 2017.
 - ^(w) lower bound given by the number of initial tokens.
 - ^(x) stated by prod in April 2017.
 - ^(y) stated by prod in April 2017.
 - ^(z) lower bound given by the number of initial tokens.
 - ^(aa) stated by prod in April 2017.
 - ^(ab) stated by prod in April 2017.
 - ^(ac) lower bound given by the number of initial tokens.
 - ^(ad) stated by prod in April 2017.
 - ^(ae) stated by prod in April 2017.
 - ^(af) lower bound given by the number of initial tokens.
 - ^(ag) lower bound given by the number of initial tokens.
 - ^(ah) lower bound given by the number of initial tokens.
 - ^(ai) lower bound given by the number of initial tokens.
 - ^(aj) lower bound given by the number of initial tokens.
 - ^(ak) lower bound given by the number of initial tokens.
 - ^(al) lower bound given by the number of initial tokens.
 - ^(am) lower bound given by the number of initial tokens.
 - ^(an) lower bound given by the number of initial tokens.
 - ^(ao) lower bound given by the number of initial tokens.
 - ^(ap) lower bound given by the number of initial tokens.
 - ^(aq) lower bound given by the number of initial tokens.
 - ^(ar) lower bound given by the number of initial tokens.
 - ^(as) lower bound given by the number of initial tokens.
 - ^(at) lower bound given by the number of initial tokens.
 - ^(au) lower bound given by the number of initial tokens.